# **SGXBounds** Memory Safety for Shielded Execution

*Dmitrii Kuvaiskii* †, Oleksii Oleksenko †, Sergei Arnautov †, Bohdan Trach †, Pramod Bhatotia \*, Pascal Felber ‡, Christof Fetzer †

**†**TU Dresden, **\*** The University of Edinburgh, **‡**University of Neuchâtel







• Security is a key barrier to adoption of cloud computing





# **Microsoft Azure**

• Security is a key barrier to adoption of cloud computing

1

• Attackers compromise **confidentiality** and **integrity** 



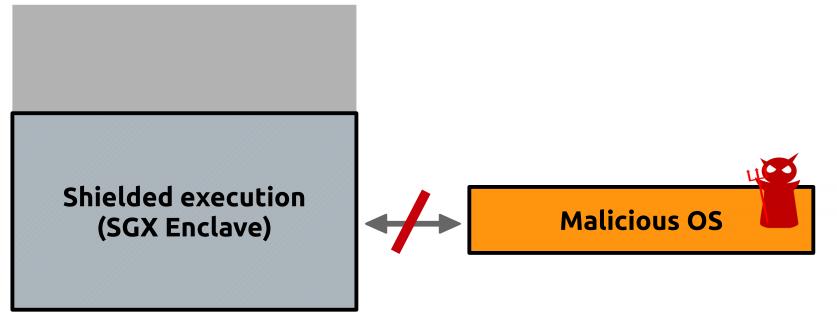
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  - → Malicious host (e.g., cloud provider)
  - → Software vulnerabilities

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Virtual Address Space

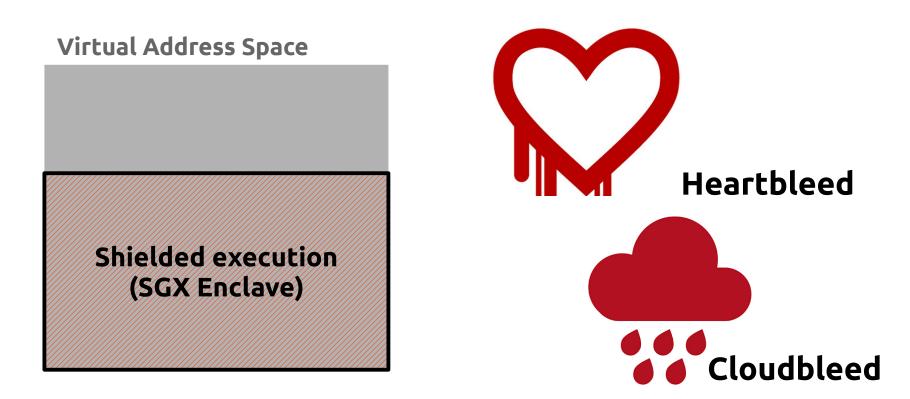


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### Protecting against Attacks

# SGX Enclave

(malicious host)

### Protecting against Attacks

# SGX Enclave (malicious host) + Memory safety (vulnerabilities)

## SGX Enclave (malicious host) H Memory safety (vulnerabilities)

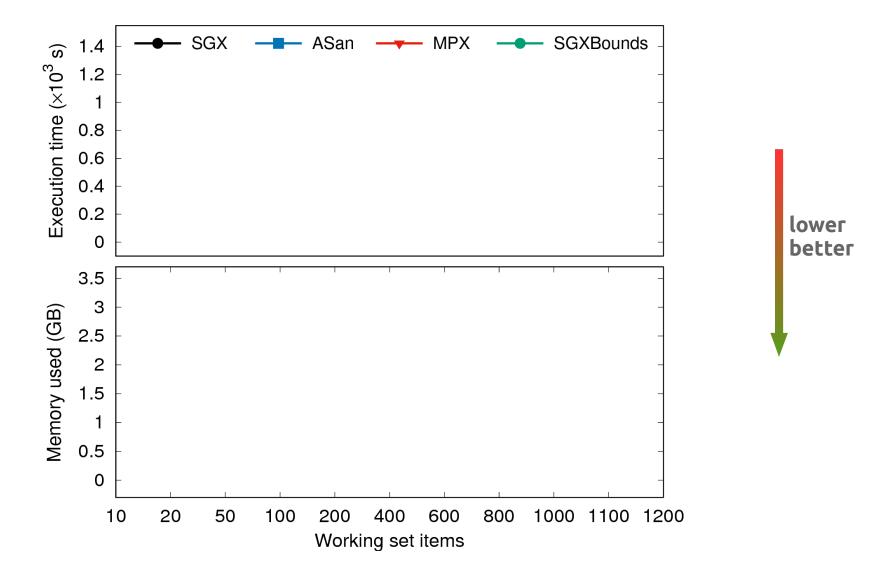
#### AddressSanitizer (software-based)

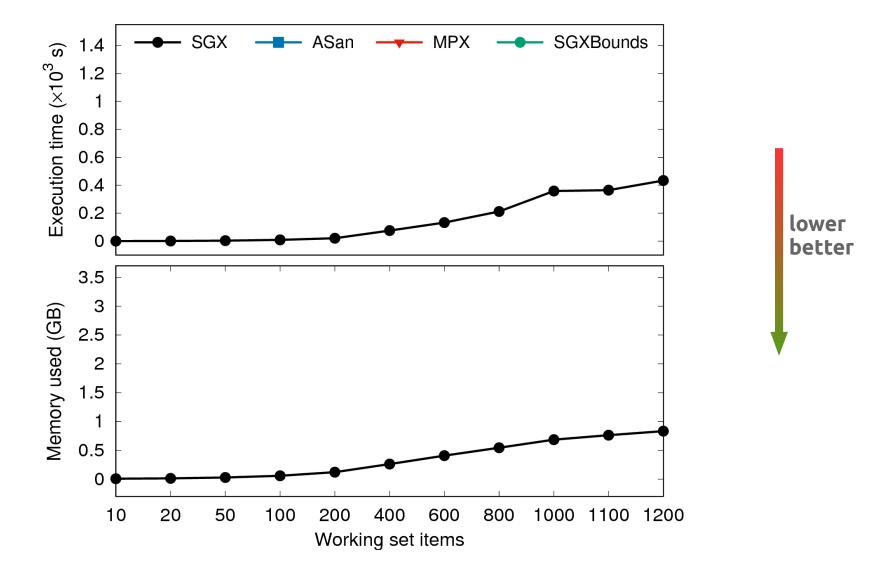
Intel MPX (hardware-based)

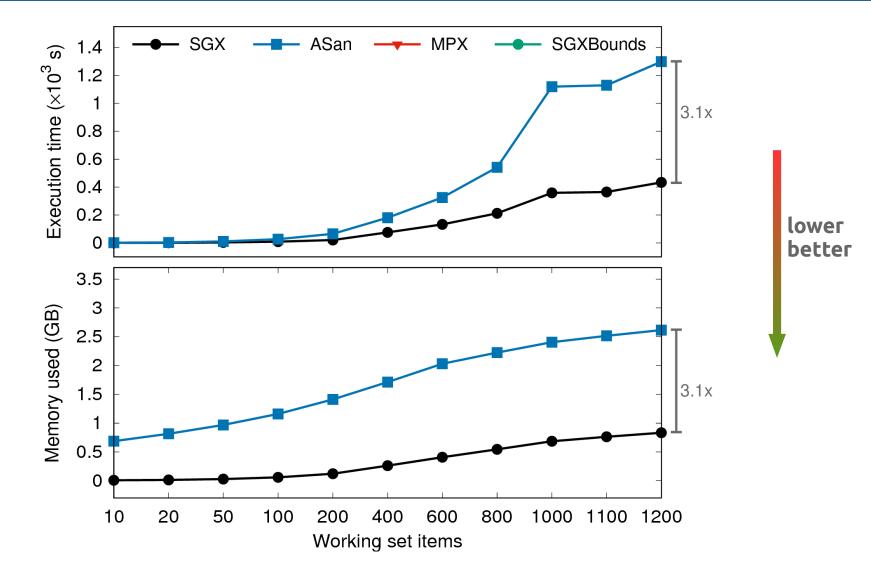
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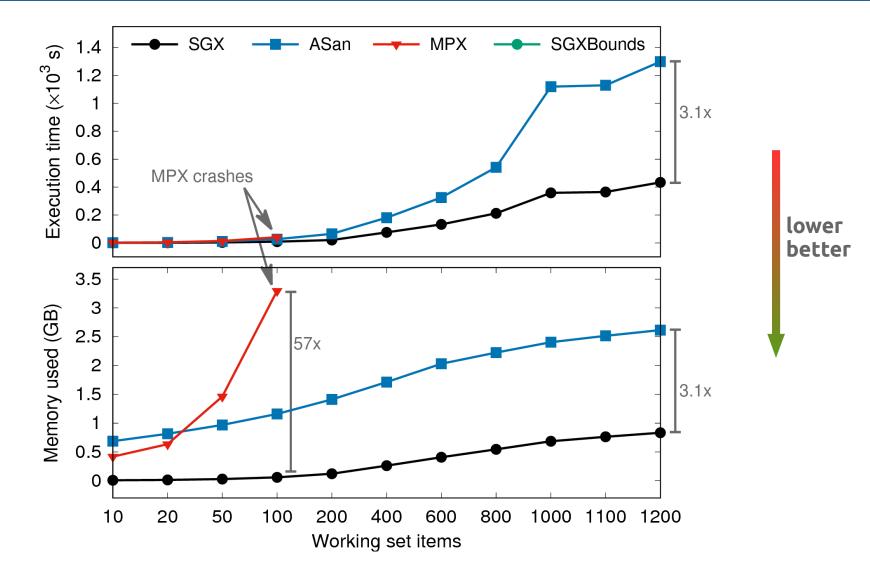
AddressSanitizer (software-based) Intel MPX (hardware-based)

# State-of-the-art memory-safety mechanisms are inefficient!



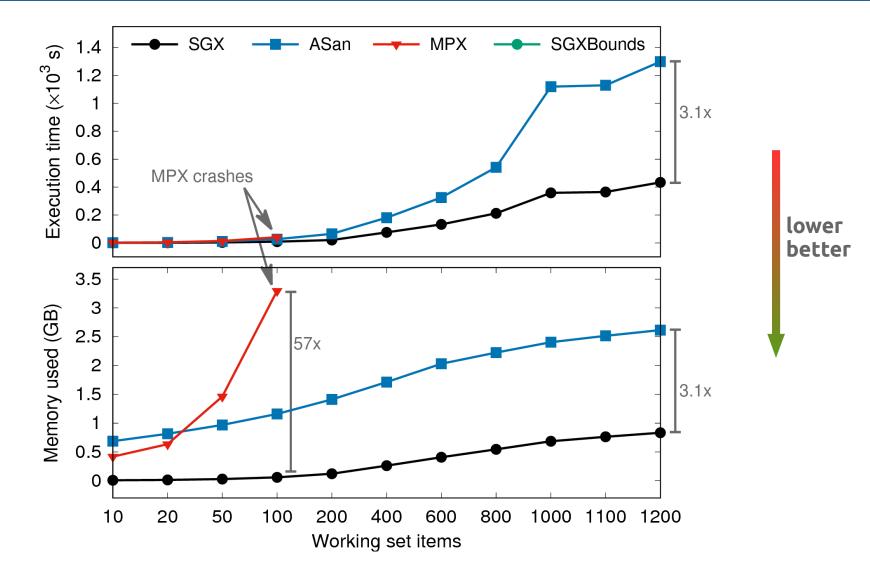


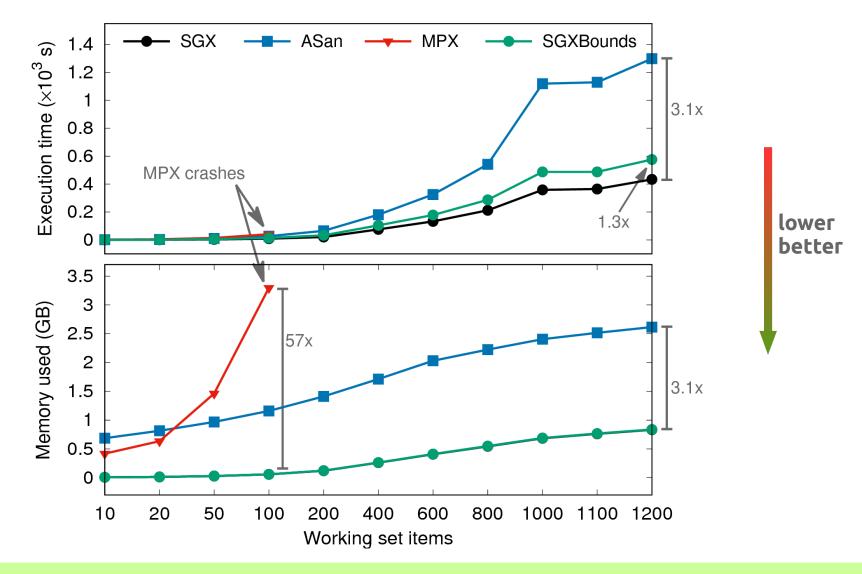




#### 

# How to make it efficient?





# SGXBounds is practical

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### - Motivation

- Constraints of SGX enclaves
- Design of SGXBounds
- Implementation of SGXBounds
- Evaluation

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### Why AddressSanitizer and Intel MPX perform poorly under SGX?

#### Virtual Address Space

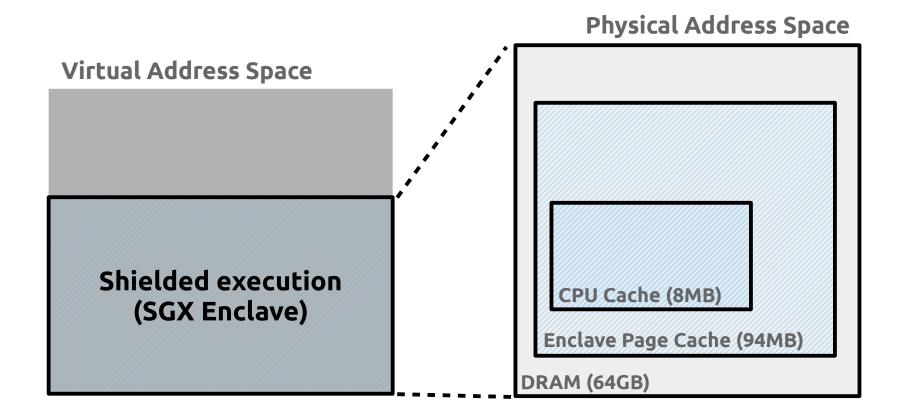
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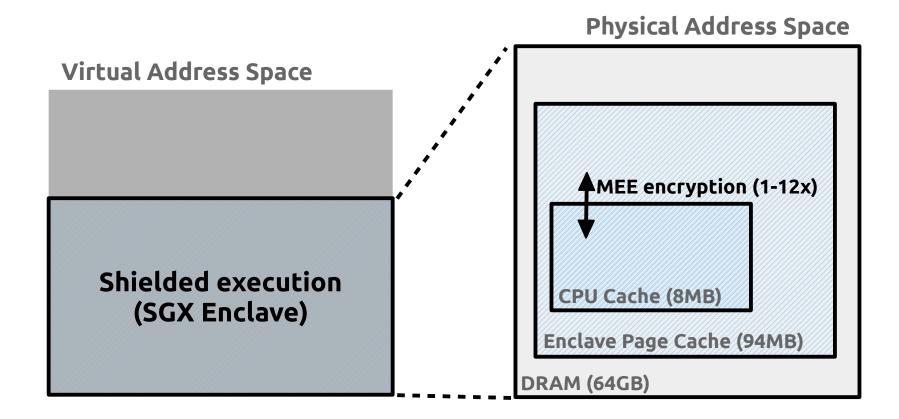
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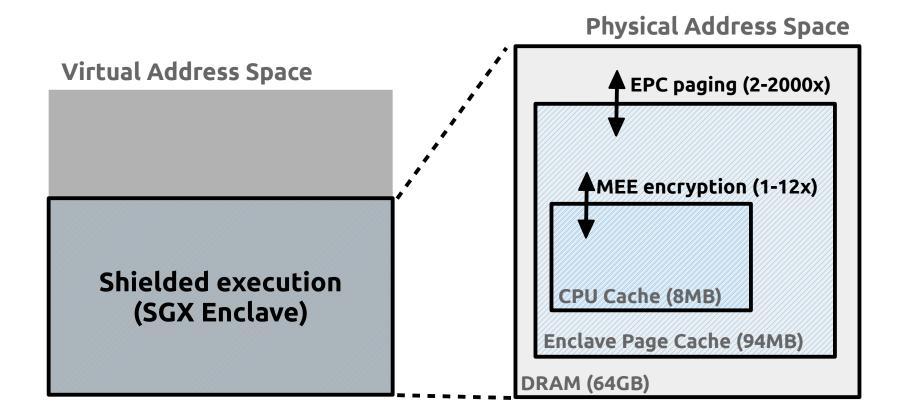
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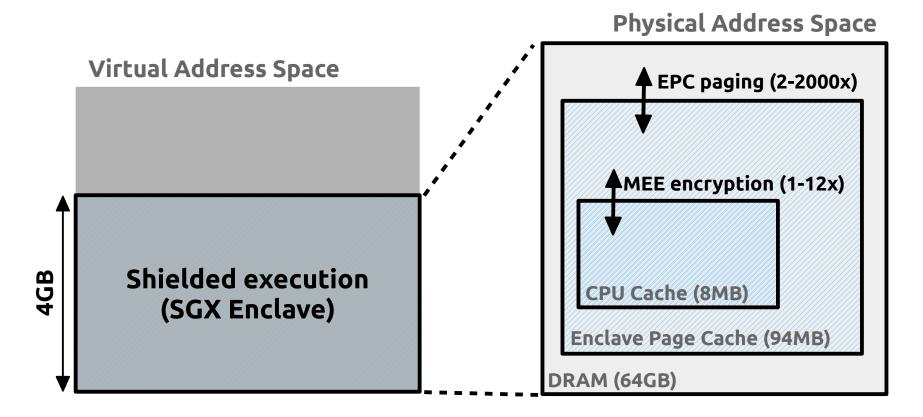


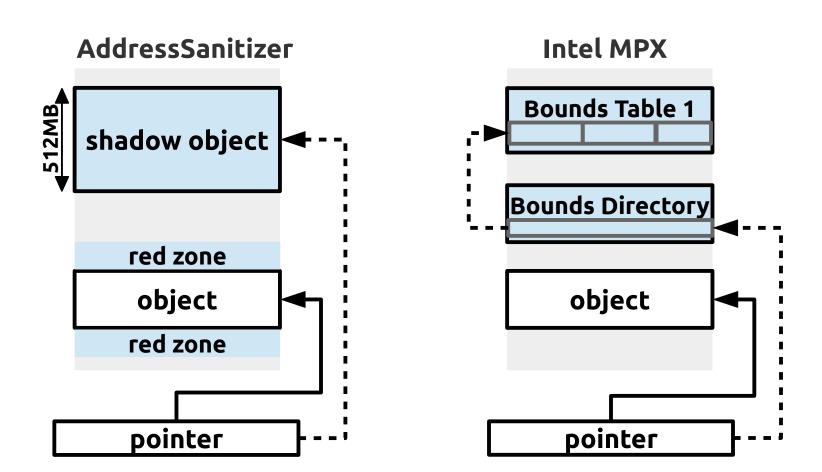
## 



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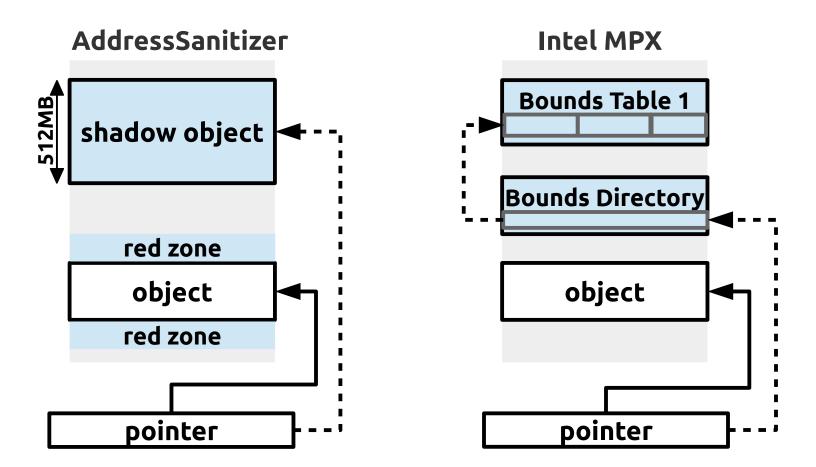
- 🙁 Increased latency of memory accesses
- 🙁 Limited enclave memory (4GB)



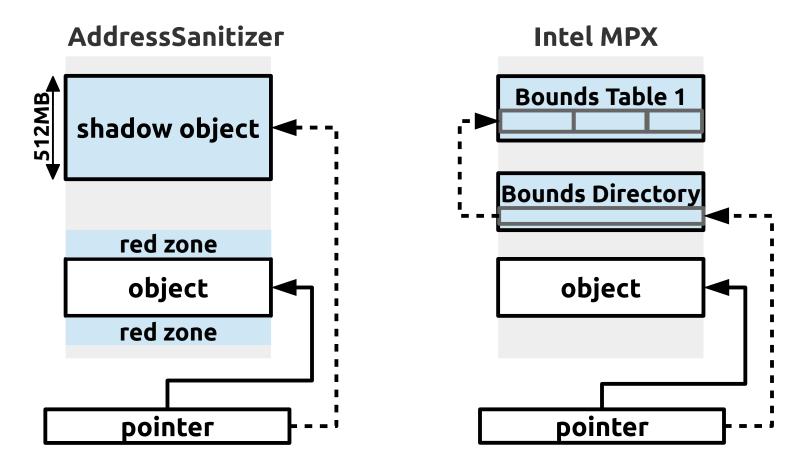


### State-of-the-Art: Metadata Layout

- 😕 Fast accesses to metadata
- 🙁 Almost endless memory



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- Almost endless memory ≠ limited enclave memory



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### - Motivation

### - Constraints of SGX enclaves

# – Design of SGXBounds

- Implementation of SGXBounds
- Evaluation

### Memory contraints of SGX dictated design of SGXBounds

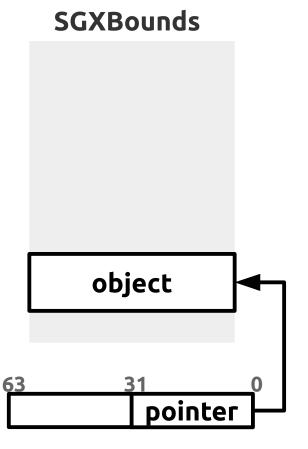
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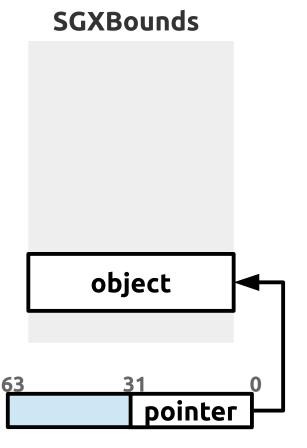
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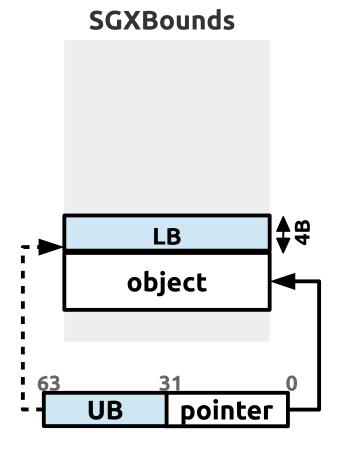
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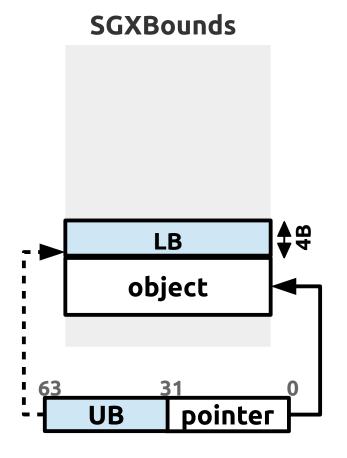
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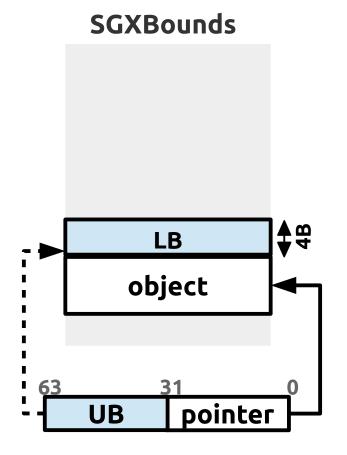


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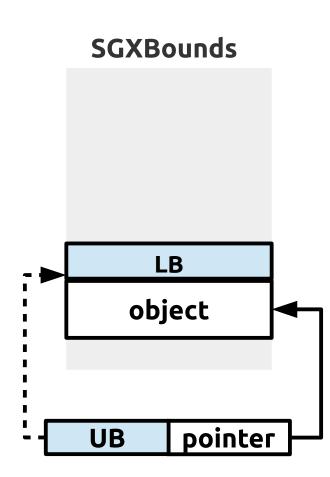
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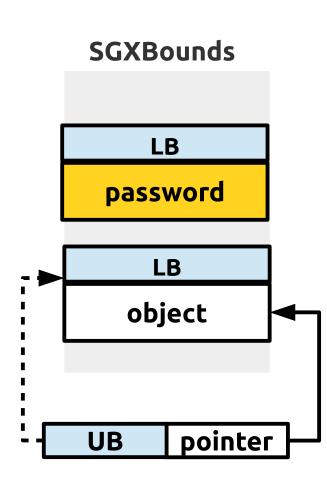


- Upper bound (UB) in pointer
- Lower bound (LB) per object
- Out-of-the-box **multithreading** (unlike MPX)

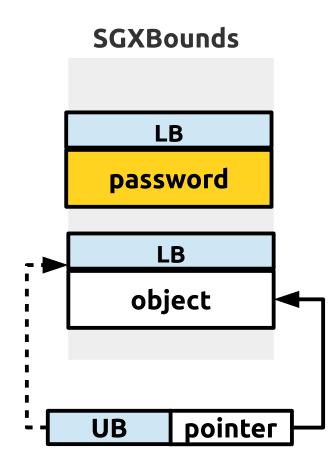
#### How SGXBounds **detects vulnerabilities** like Heartbleed?



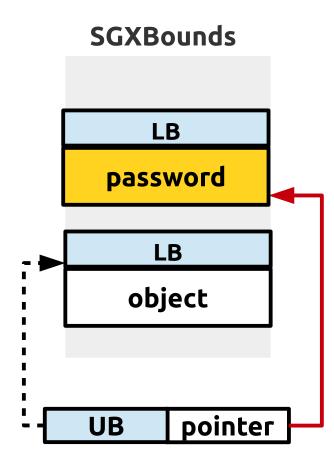
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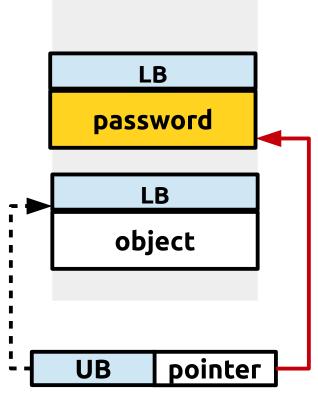


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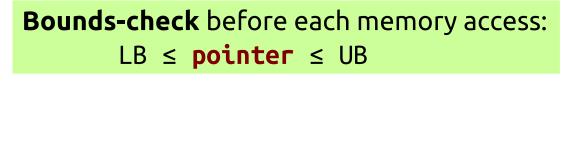


How SGXBounds **detects vulnerabilities** like Heartbleed?

- Oata leak through write(socket, pointer, objlen)
- Protect using efficient bounds checks

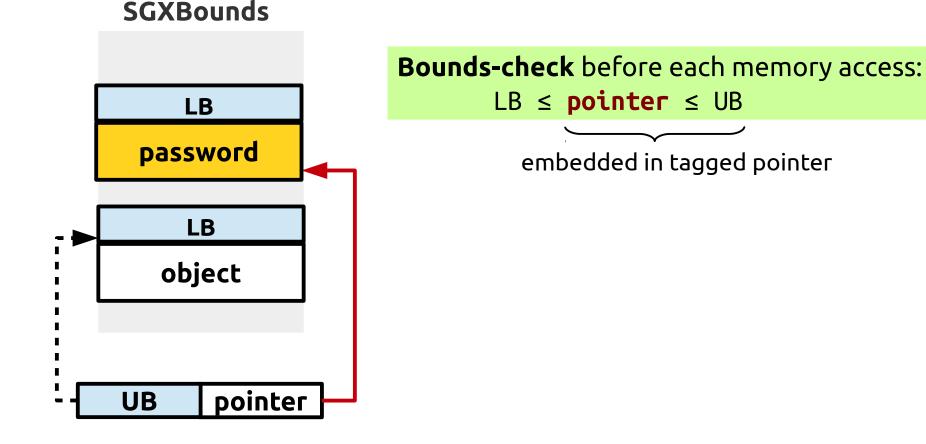


SGXBounds



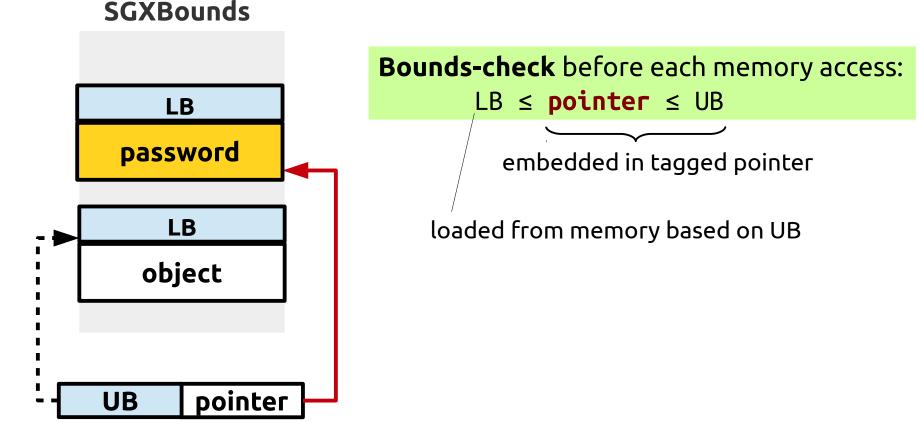
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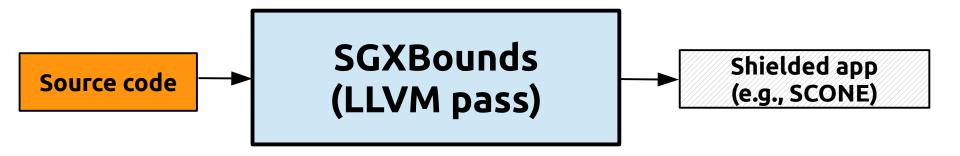
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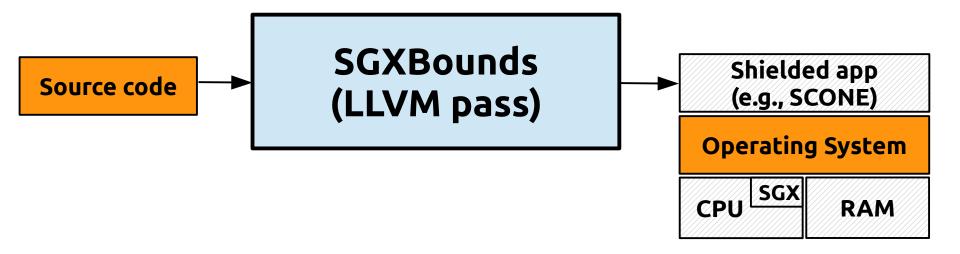
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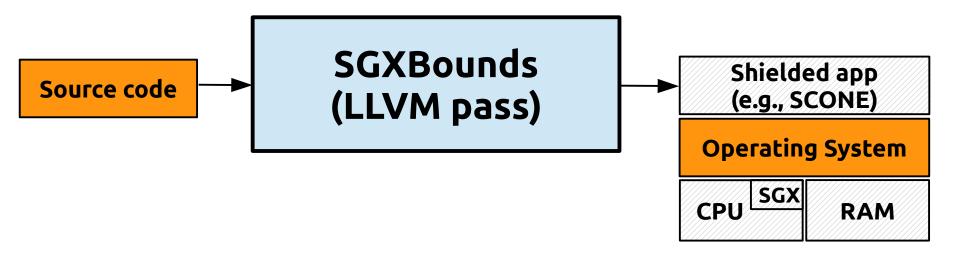


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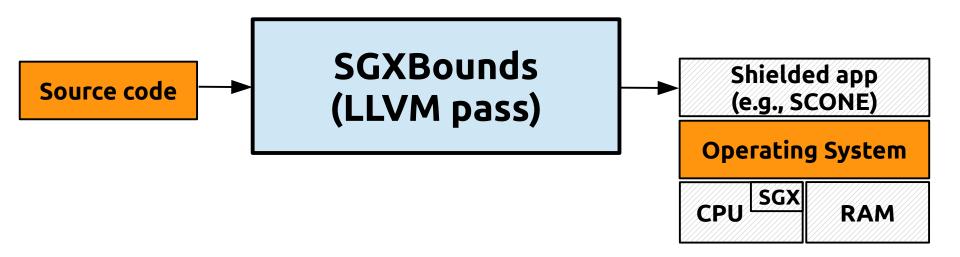
# SGXBounds (LLVM pass)





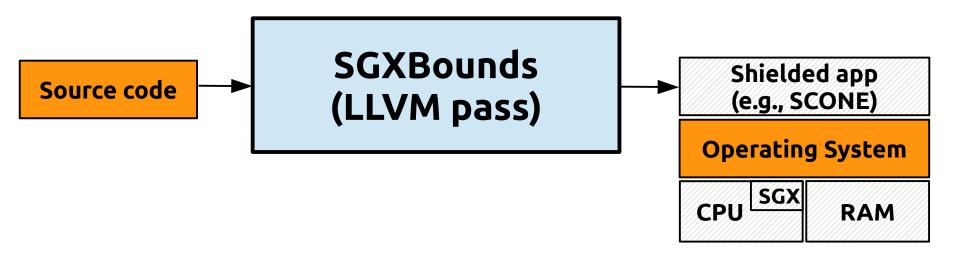


#### Advanced features:



Advanced features:

- → Tolerating errors with boundless memory
- → Metadata management support
- → Compile-time **optimizations**



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#### See paper for details

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  - → Benchmark suites
  - → Case studies
  - ➡ Security

ASan	ΜΡΧ	SGXBounds

	ASan	ΜΡΧ	SGXBounds
Phoenix	1.41	2.27	1.13

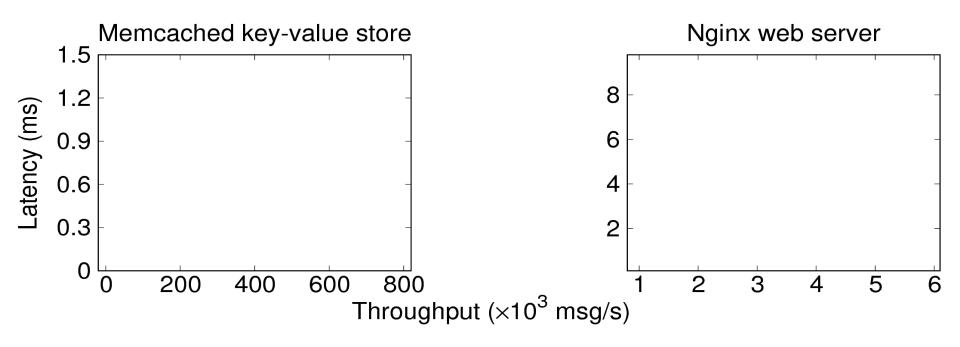
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PARSEC	1.60	1.43*	1.20

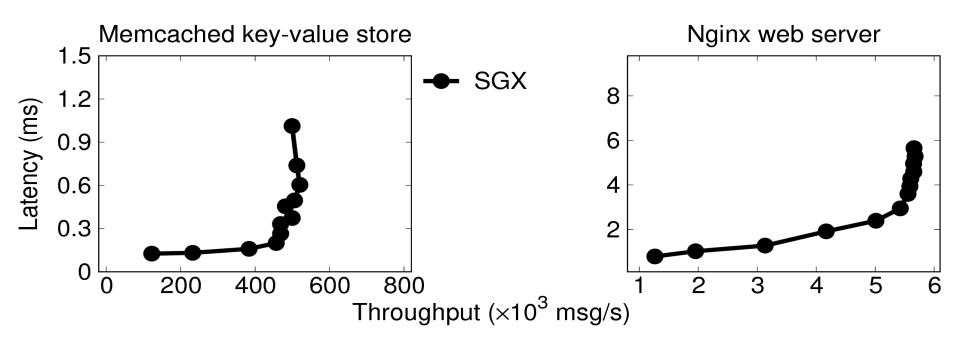
\* some programs failed due to insufficient memory

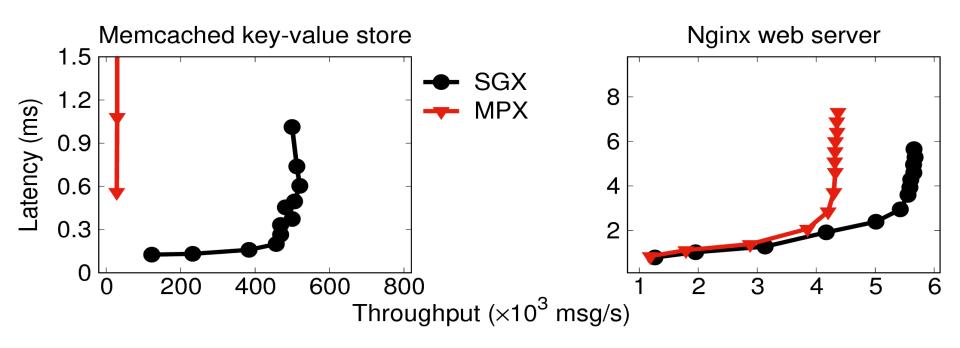
	ASan	ΜΡΧ	SGXBounds
Phoenix	1.41	2.27	1.13
PARSEC	1.60	1.43*	1.20
SPEC	1.76	1.52*	1.41

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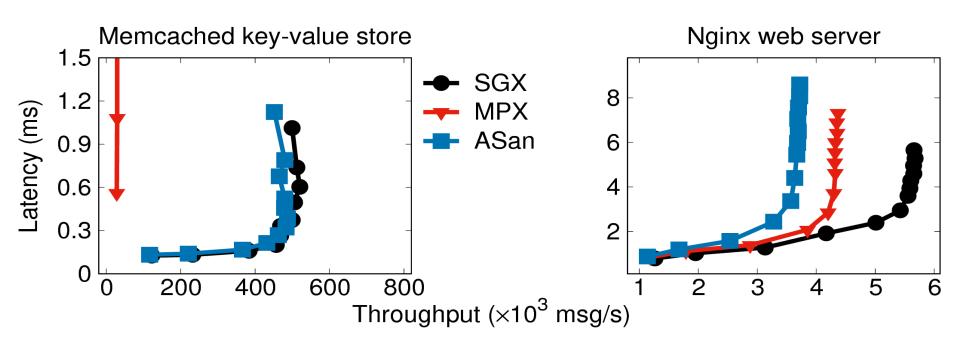
# **Case Studies**



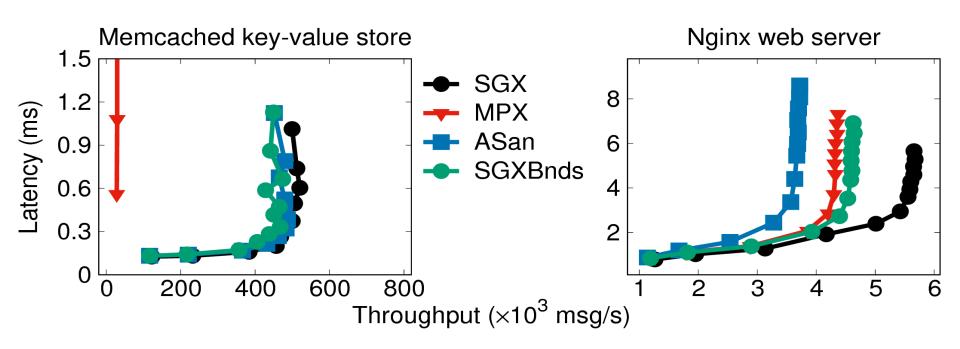




😕 MPX: EPC thrashing on Memcached



# MPX: EPC thrashing on Memcached ASan: metadata overload on Nginx



MPX: EPC thrashing on Memcached
 ASan: metadata overload on Nginx
 SGXBounds: no corner cases

### Security guarantees

# 😳 **RIPE** synthetic benchmark:

→ Similar guarantees as ASan and MPX

# **C RIPE** synthetic benchmark:

#### → Similar guarantees as ASan and MPX

# Real-world vulnerabilities detected and tolerated:

- → Memcached denial-of-service
- → Nginx stack buffer overflow
- → Apache Heartbleed

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## Conclusion

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 Use shielded execution with Intel SGX

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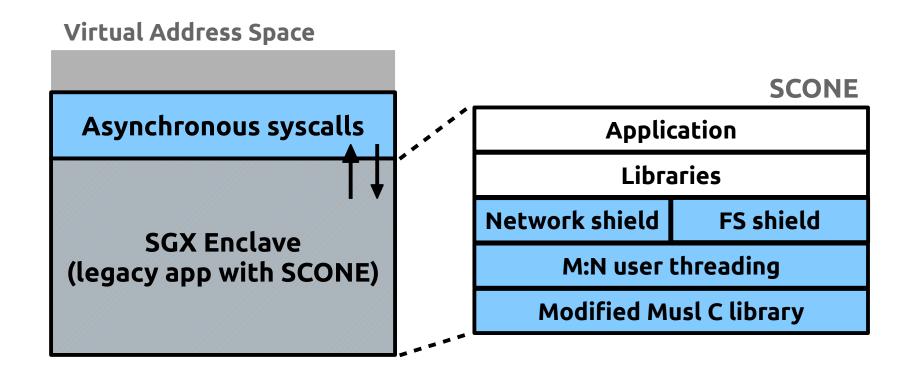
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# Thank you!

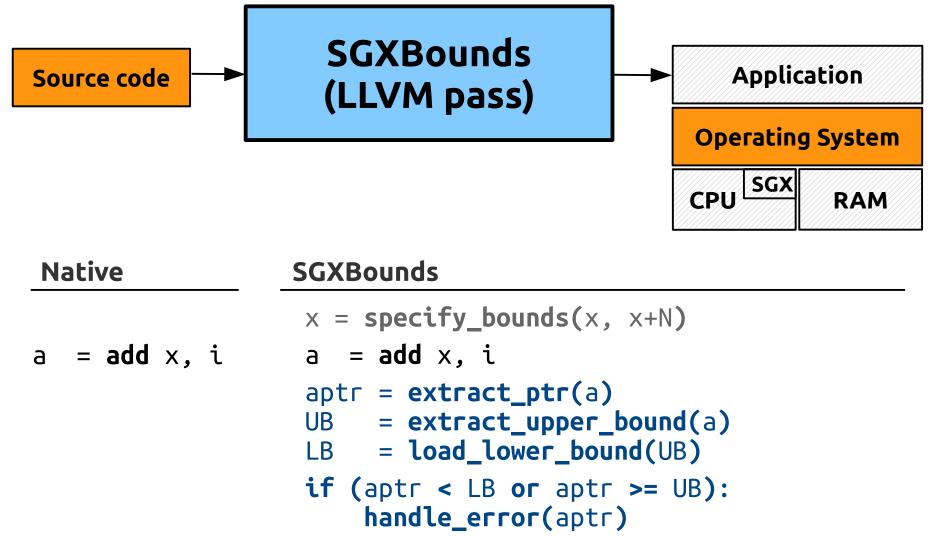
dmitrii.kuvaiskii@tu-dresden.de https://github.com/tudinfse/sgxbounds

## Backup slides



-V. Costan, S. Devadas. "Intel SGX Explained". IACR Cryptology ePrint Archive '16 -S. Arnautov et al. "SCONE: Secure linux containers with Intel SGX". OSDI'16 Eurosys 2017

# **SGXBounds: Implementation**

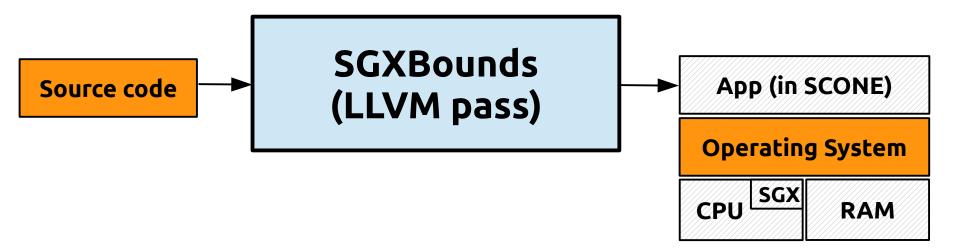


store 42, a store 42, a

	Perf	Mem	Comments
Intel MPX	146%	116%	FP/FN for multithreaded
AddressSanitizer	38%	292%	-
BaggyBounds <sup>1</sup>	70%	12%	Not publicly available
Low-Fat Pointers <sup>2</sup>	54%	12%	Not publicly available
SGXBounds	55%	0%	(this work)

<sup>1</sup> P. Akritidis et al. "Baggy Bounds Checking: An efficient and backwards-compatible defense against out-of-bounds errors". Usenix Security'09
 <sup>2</sup> G. Duck et al. "Stack Bounds Protection with Low Fat Pointers". NDSS'17

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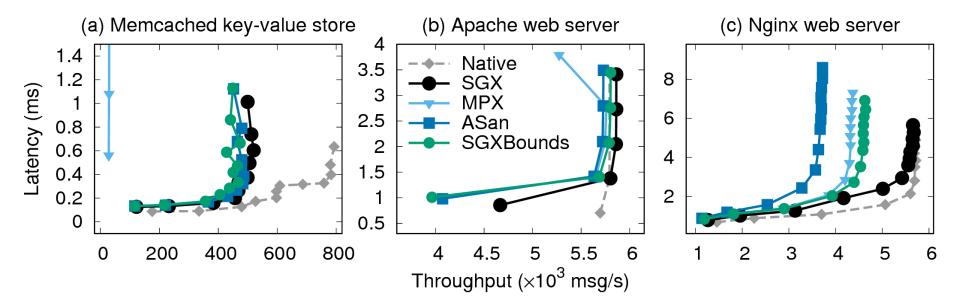
#### Instrumentation:

data:lower bound metadata after each allocated objectpointers:upper bound metadata in each data pointercode:bounds-check before each memory access

# Security guarantees

D detected?T tolerated?

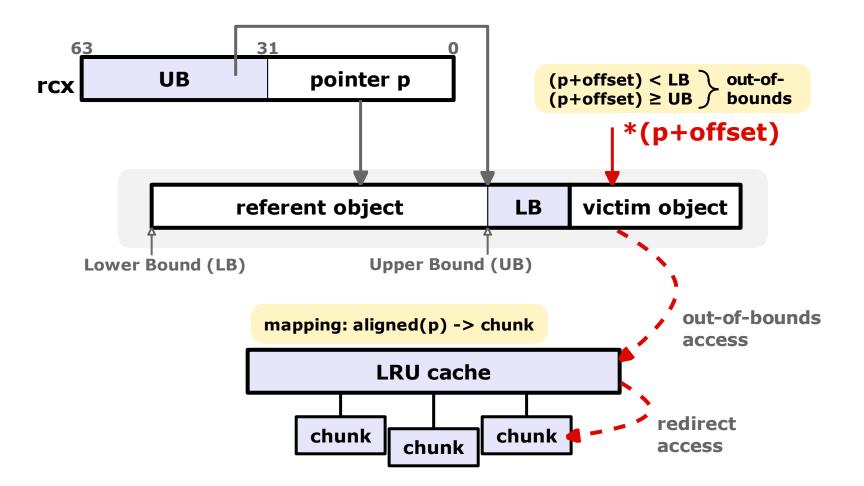
	MPX	ASan	SGXBounds	
RIPE benchmark	2/16	8/16	8/16	
Memcached CVE-2011-4971	D (T)	D (T)	D (T)	
Nginx CVE-2013-2028	D (T)	D (T)	D (T)	
Apache Heartbleed	D (T)	D (T)	D (T)	



	CF	DO	IL
Control Flow Integrity [27, 39, 52, 84]	~	×	×
Code Pointer Integrity [46]		×	×
Address Space Randomization [45, 48, 50, 68, 70]	✔*	×	×
Data Integrity [16]	~	~	×
Data Flow Integrity [29]	~	~	×
Software Fault Isolation [39, 79]	V	V	V
Data Space Randomization [24, 28]	*	*	*
Memory safety [9, 17, 20, 26, 35, 55, 58, 69]	~	~	~

\*SGX enclaves do not provide sufficient bits of entropy in random offsets/masks

**CF** – control flow hijack, **DO** – data-only attack, **IL** – information leak



<sup>1</sup> M. Rinard et al. "A dynamic technique for eliminating buffer overflow vulnerabilities (and other memory errors)". ACSAC'04

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## **SGXBounds: Outside of Enclaves**

